



Analysis and Design of Algorithms

Lecture 3



Analysis of Algorithms II

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Maximum Pairwise Product

□ Given a sequence of non-negative integers a_0, \dots, a_{n-1} , find the **maximum pairwise product**, that is, the largest integer that can be obtained by multiplying two different elements from the sequence (or, more formally, $\max a_i a_j$ where $0 \leq i \neq j \leq n-1$). Different elements here mean a_i and a_j with $i \neq j$ (it can be the case that $a_i = a_j$).

Maximum Pairwise Product

□ Constraints $2 \leq n \leq 2 * 10^5$ & $0 \leq a_0, \dots, a_{n-1} \leq 10^5$.

Maximum Pairwise Product

□ Sample 1

❖ Input: 1 2 3

❖ Output: 6

□ Sample 2

❖ Input: 7 5 14 2 8 8 10 1 2 3

❖ Output: 140

Maximum Pairwise Product

□ Sample 3

❖ Input: 4 6 2 6 1

❖ Output: 36

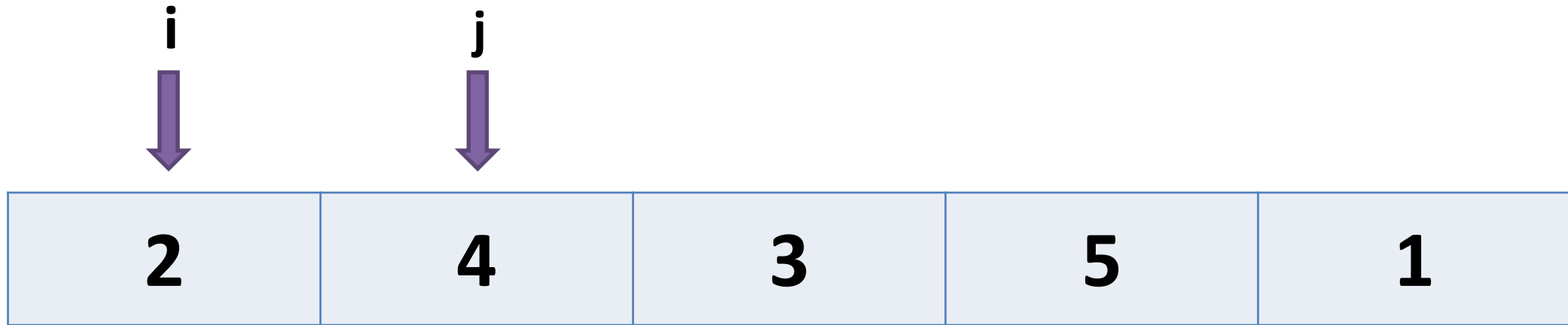
Maximum Pairwise Product

□ Assume the following array

| | | | | |
|----------|----------|----------|----------|----------|
| 2 | 4 | 3 | 5 | 1 |
|----------|----------|----------|----------|----------|

Maximum Pairwise Product

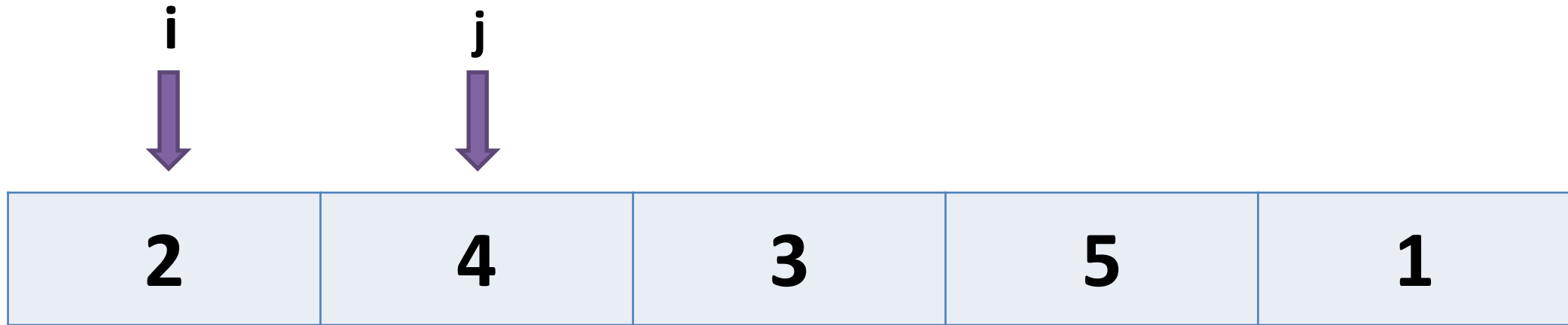
□ Assume the following array



Result=0

Maximum Pairwise Product

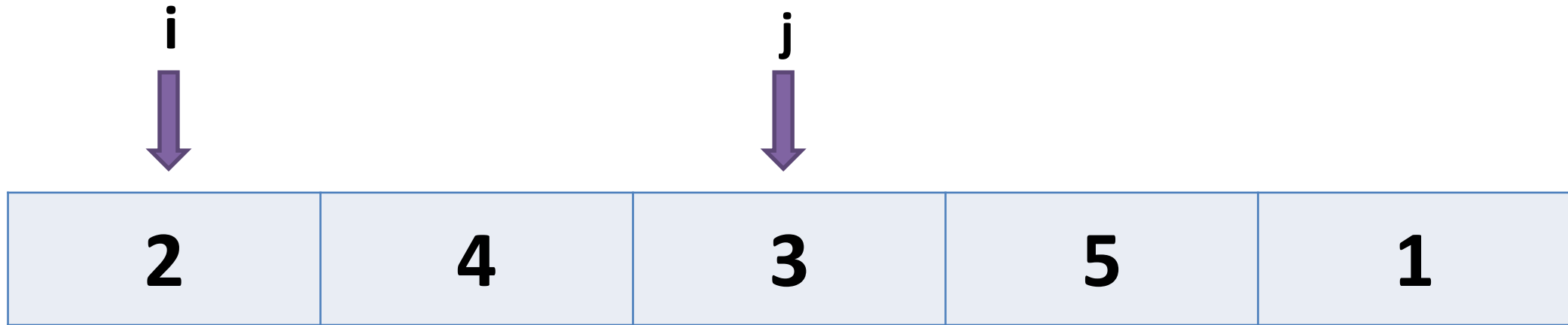
□ Assume the following array



If $a[i] * a[j] > \text{result}$
 $\text{result} = a[i] * a[j] = 8$

Maximum Pairwise Product

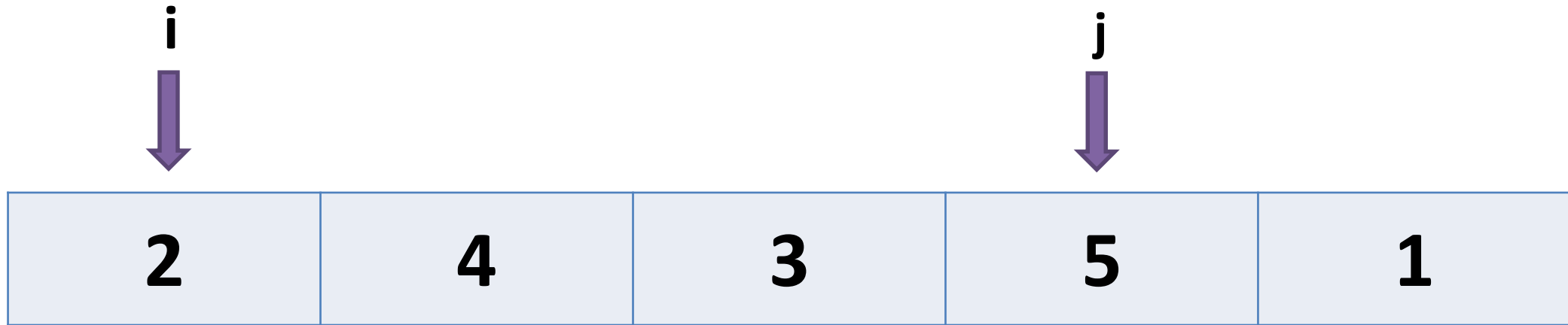
□ Assume the following array



If $a[i] * a[j] > \text{result}$
 $\text{result} = 8$

Maximum Pairwise Product

□ Assume the following array



If $a[i] * a[j] > \text{result}$
 $\text{result} = a[i] * a[j] = 10$

Maximum Pairwise Product

- Naive algorithm

```
def max_pairwise_product(a):  
    result = 0  
    for i in range(0, len(a)):  
        for j in range(i+1, len(a)):  
            if a[i]*a[j] > result:  
                result = a[i]*a[j]  
    return result
```

Maximum Pairwise Product

□ Python Code:

```
a=[2,6,4,5,2]  
print(max_pairwise_product(a))
```

30

Maximum Pairwise Product

□ Time complexity $O(n^2)$

```
def max_pairwise_product(a):  
    result = 0  
    for i in range(0, len(a)):  
        for j in range(i+1, len(a)):  
            if a[i]*a[j] > result:  
                result = a[i]*a[j]  
    return result
```

Maximum Pairwise Product

□ we need a faster algorithm. This is because our program performs about n^2 steps on a sequence of length n . For the maximal possible value $n=200,000 = 2 \cdot 10^5$, the number of steps is $40,000,000,000 = 4 \cdot 10^{10}$. This is too much. Recall that modern machines can perform roughly 10^9 basic operations per second

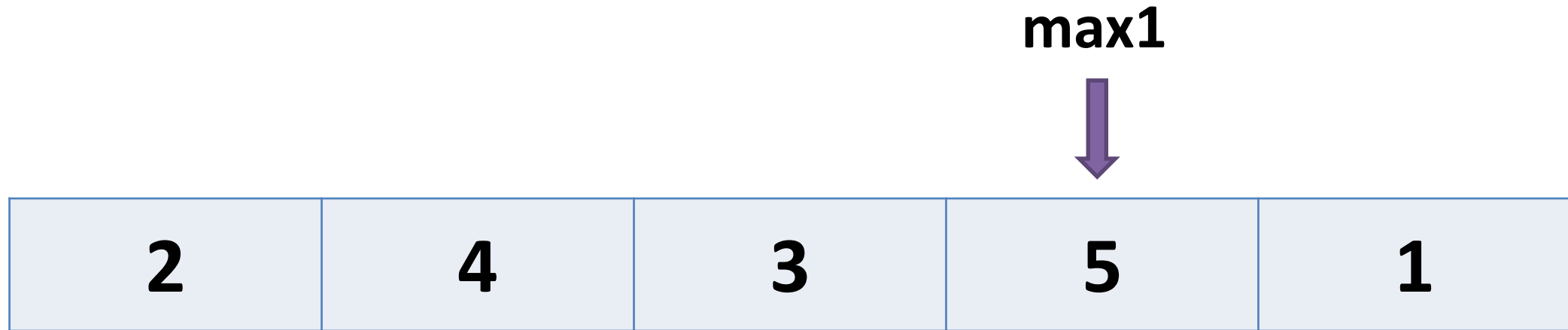
Maximum Pairwise Product

□ Assume the following array

| | | | | |
|----------|----------|----------|----------|----------|
| 2 | 4 | 3 | 5 | 1 |
|----------|----------|----------|----------|----------|

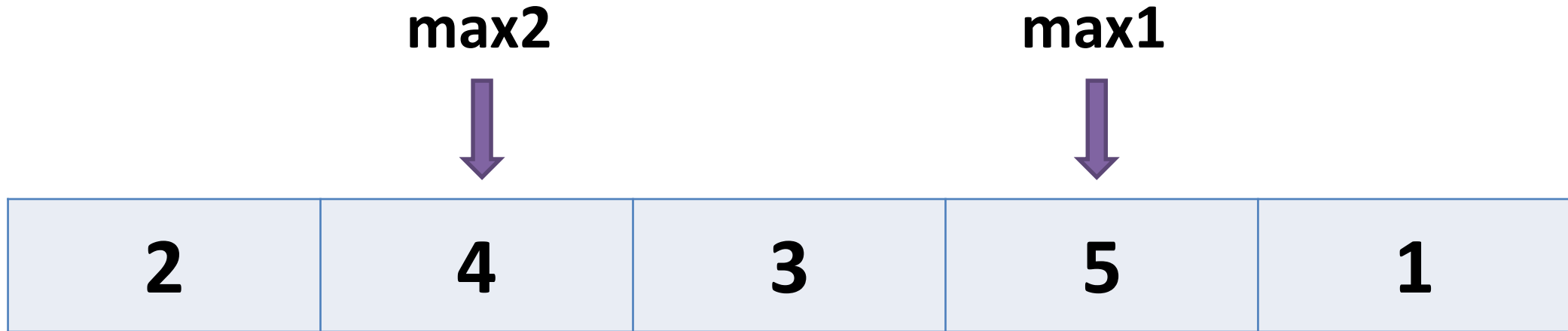
Maximum Pairwise Product

- Find maximum number 1



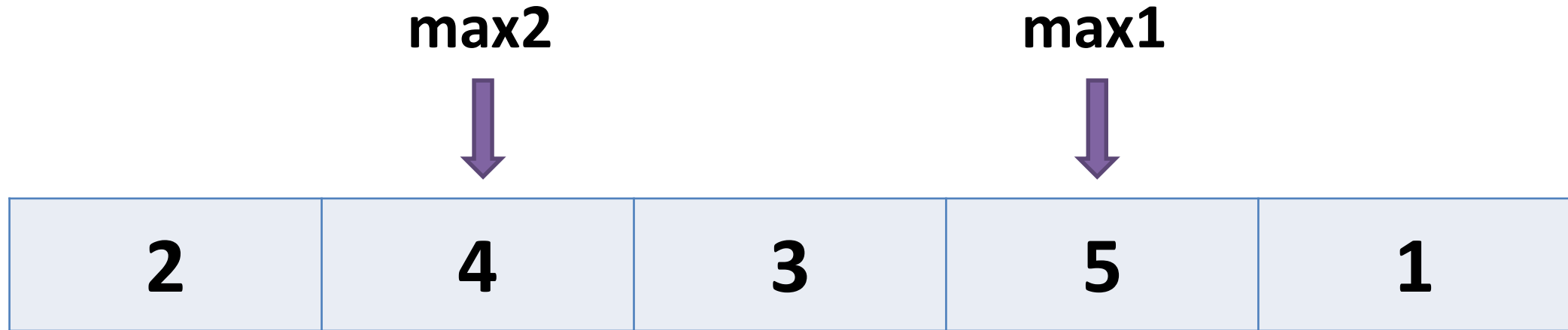
Maximum Pairwise Product

- Find maximum number2 but not maximum number1



Maximum Pairwise Product

- Find maximum number2 but not maximum number1



return max1*max2

Maximum Pairwise Product

□ Efficient
algorithm

```
def max_pairwise_product_fast(numbers):  
    max1 = -1  
    index1 = None  
    max2 = -1  
    for element in range(len(numbers)):  
        if numbers[element] >= max1:  
            max1 = numbers[element]  
            index1 = element  
    for element in range(len(numbers)):  
        if numbers[element] >= max2:  
            if element != index1:  
                max2 = numbers[element]  
    return max1 * max2
```

Maximum Pairwise Product

- Efficient algorithm

```
numbers=[2,6,4,5,2]  
print(max_pairwise_product_fast(numbers))
```

30

Maximum Pairwise Product

□ Time complexity $O(n)$

```
def max_pairwise_product_fast(numbers):  
    max1 = -1  
    index1 = None  
    max2 = -1  
    for element in range(len(numbers)):  
        if numbers[element] >= max1:  
            max1 = numbers[element]  
            index1 = element  
    for element in range(len(numbers)):  
        if numbers[element] >= max2:  
            if element != index1:  
                max2 = numbers[element]  
    return max1 * max2
```

Fibonacci

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

□ *Definition:*

$$\square F_n = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ F_{n-2} + F_{n-1} & n > 1 \end{cases}$$

□ Examples:

$$\diamond F_8 = 21$$

$$\diamond F_{20} = 6765$$

$$\diamond F_{50} = 12586269025$$

$$\diamond F_{100} = 354224848179261915075$$

Fibonacci

□ Examples:

❖ $F_{500} =$

1394232245616978801397243828

7040728395007025658769730726

4108962948325571622863290691

557658876222521294125

Fibonacci

- Naive algorithm

```
def fib(n):  
    if (n <= 1):  
        return n  
    else:  
        return fib(n-1) + fib(n-2)
```

Fibonacci

- Naive algorithm

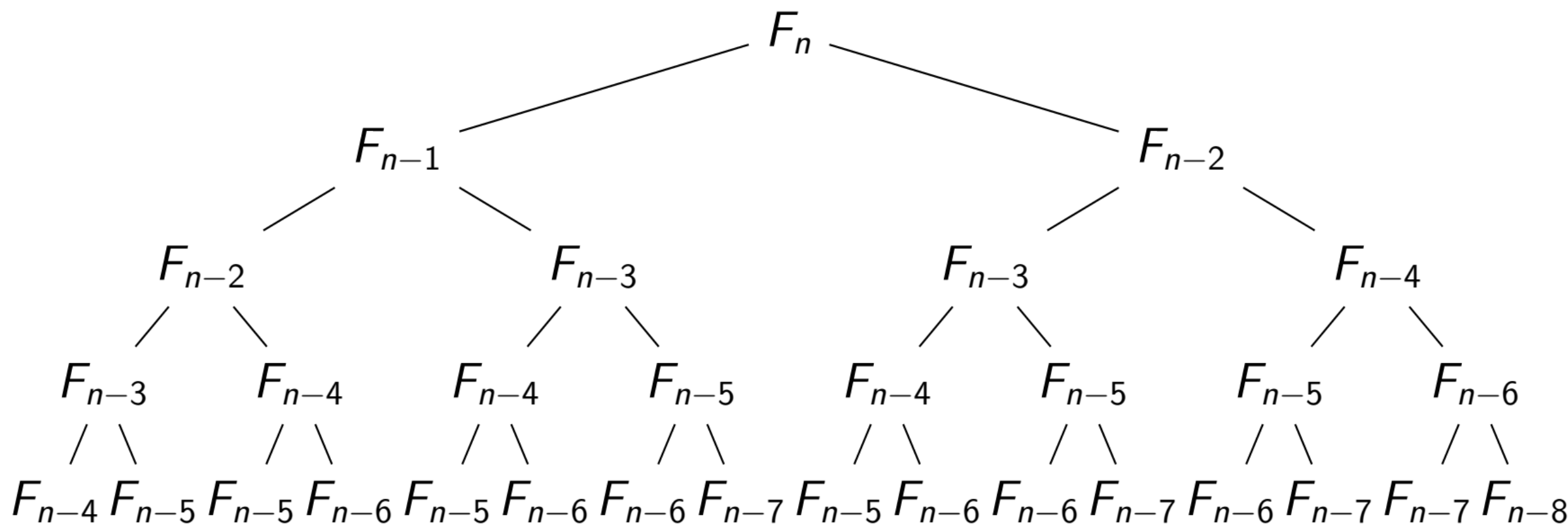
```
print(fib(20))
```

6765

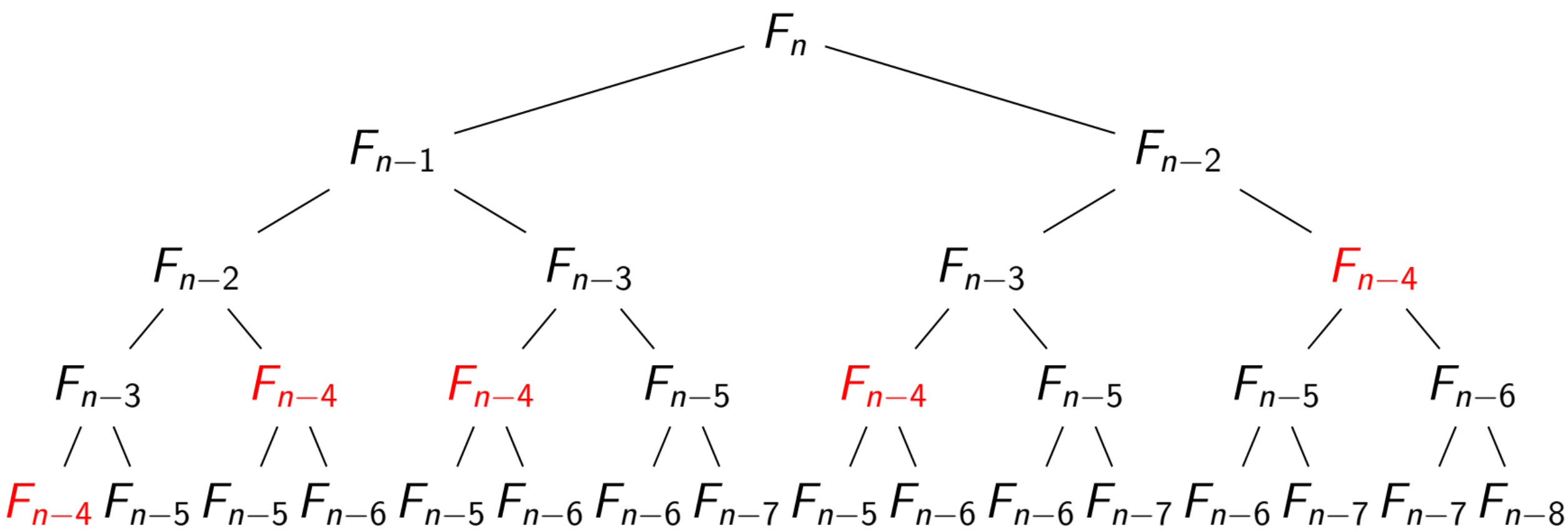
```
print(fib(100))
```

Very Long Time why????

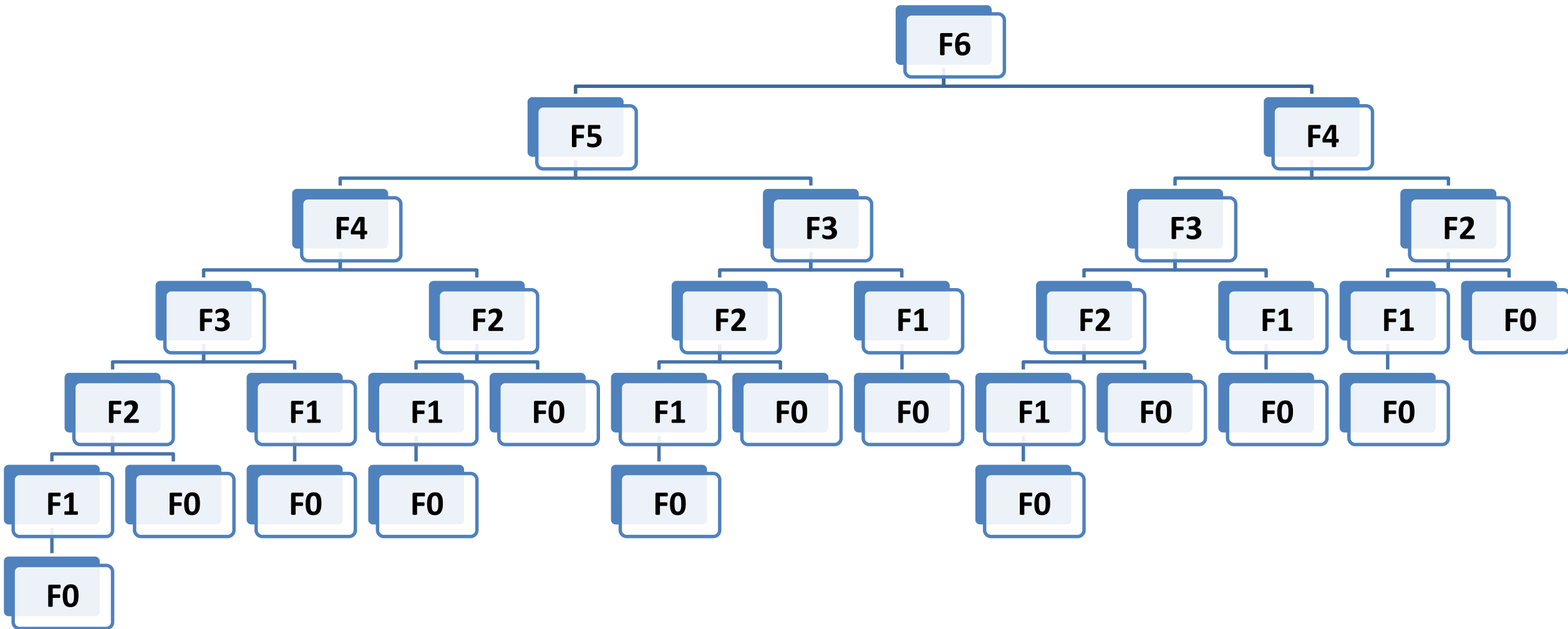
Fibonacci



Fibonacci



Fibonacci

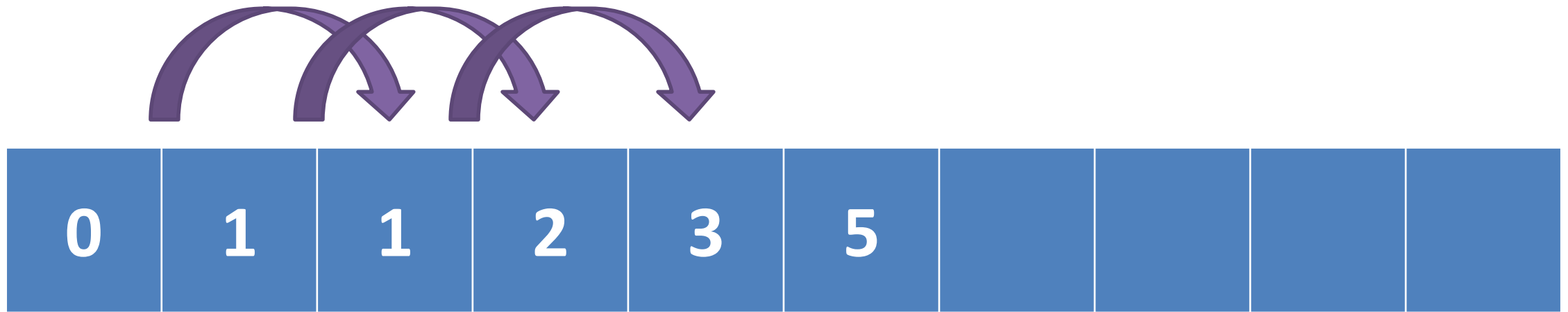


❑ Fib algorithm is very slow because of recursion

❑ Time complexity = $O(2^n)$

Fibonacci

- Efficient algorithm



Create array then insert fibonacci

Fibonacci

□ Efficient algorithm

```
def fib_fast(n):  
    if (n <= 1):  
        return n  
    else:  
        numbers = [0, 1]  
        for i in range(n-1):  
            numbers.append(numbers[-1]+numbers[-2])  
        return numbers[-1]
```

Fibonacci

- ❑ Efficient algorithm

```
print(fib(20))
```

6765

```
print(fib_fast(100))
```

354224848179261915075

short Time why????

- ❑ Fib_Fast algorithm is fast because of
loop + array
- ❑ Time complexity = $O(n^2)$

Fibonacci

❑ Efficient algorithm

➤ Try

```
print(fib_faster(1000000))
```

Very long Time why????

- Advanced algorithm
- ❖ No array
- ❖ Need two variable + Loop

Fibonacci

- ❑ Advanced algorithm
- ❑ Compute F_6
- ❑ $a=0, b=1$

| a | b |
|----------|----------|
| 0 | 1 |

Fibonacci

- ❑ Advanced algorithm
- ❑ Compute F_6
- ❑ $a=b, b=a+b$



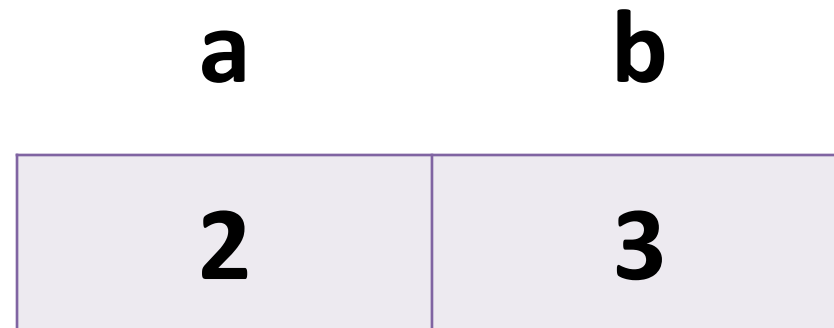
Fibonacci

- ❑ Advanced algorithm
- ❑ Compute F_6
- ❑ $a=b, b=a+b$



Fibonacci

- ❑ Advanced algorithm
- ❑ Compute F_6
- ❑ $a=b, b=a+b$



Fibonacci

- ❑ Advanced algorithm
- ❑ Compute F_6
- ❑ $a=b, b=a+b$

| a | b |
|----------|----------|
| 3 | 5 |

Fibonacci

- ❑ Advanced algorithm
- ❑ Compute F_6
- ❑ $a=b, b=a+b$

| a | b |
|----------|----------|
| 5 | 8 |

Fibonacci

- ❑ Advanced algorithm
- ❑ Compute $F_6=8$

| a | b |
|----------|----------|
| 5 | 8 |

Fibonacci

□ Advanced algorithm

```
def fib_faster(n):  
    if (n <= 1):  
        return n  
    else:  
        a, b = 0, 1  
        for i in range(n-1):  
            b, a = b + a, b  
        return b
```


Fibonacci

- ❑ Advanced algorithm

```
print(fib(20))
```

6765

```
print(fib_fast(100))
```

354224848179261915075

Very short Time why????

- ❑ Fib_Faster algorithm is faster because of loop + two variables
- ❑ Time complexity = $O(n)$

Fibonacci

□ Advanced algorithm

➤ Try

```
print(fib_faster(1000000))
```

Short Time why????

Greatest Common Divisors

Greatest Common Divisors

□ In mathematics, the greatest common divisor (gcd) of two or more integers, which are not all zero, is the largest positive integer that divides each of the integers.

Greatest Common Divisors

- Input Integers $a, b \geq 0$
- Output $\text{gcd}(a, b)$

Greatest Common Divisors

- ❑ What is the greatest common divisor of 54 and 24?
- ❑ The divisors of 54 are: 1,2,3,6,9,18,27,54
- ❑ Similarly, the divisors of 24 are: 1,2,3,4,6,8,12,24
- ❑ The numbers that these two lists share in common are the common divisors of 54 and 24: 1,2,3,6
- ❑ The greatest of these is 6. That is, the greatest common divisor of 54 and 24. $\text{gcd}(54,24)=6$

Greatest Common Divisors

□ Naive algorithm

```
def gcd(a,b):  
    if a>b:  
        i=a  
    else:  
        i=b  
    while i>=1:  
        if a%i==0 and b%i==0:  
            break  
        i=i-1  
    return i
```


Greatest Common Divisors

```
print(gcd(54, 24))
```

6

```
print(gcd(3918848, 1653264))
```

61232

- ❑ gcd algorithm is slow because of loop
- ❑ Time complexity = $O(n)$
- ❑ n depend on a, b

Greatest Common Divisors

- Efficient algorithm

```
def gcd_fast(a, b):  
    if a%b == 0:  
        return b  
    return gcd_fast(b, a%b)
```

Greatest Common Divisors

- ❑ Efficient algorithm

```
print(gcd_fast(54, 24))
```

6

```
print(gcd_fast(3918848, 1653264))
```

61232

Greatest Common Divisors

- Efficient algorithm

```
gcd_fast((3918848, 1653264))
```

```
gcd_fast((1653264, 612320))
```

```
gcd_fast((612320, 428624))
```

```
gcd_fast((428624, 183696))
```

```
gcd_fast((183696, 61232))
```

```
return 61232
```

Greatest Common Divisors

- ❑ Efficient algorithm
- ❑ Take 5 steps to solve `gcd_fast((3918848, 1653264))`
- ❑ Time complexity = $O(\log(n))$
- ❑ n depend on a, b

Summary

- ❑ Naive algorithm is too slow.
- ❑ The Efficient algorithm is much better.
- ❑ Finding the correct algorithm requires knowing something interesting about the problem.

Contact Me



**THANKS FOR
YOUR TIME**

